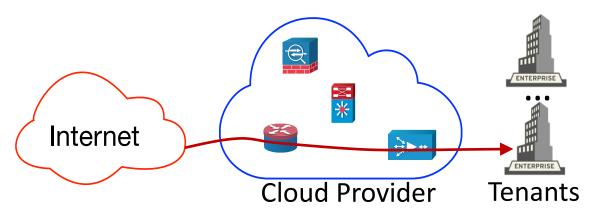
VNF Chain Allocation and Management at Datacenter Scale



Nodir Kodirov

Sam Bayless, Fabian Ruffy, Swati Goswami Ivan Beschastnikh, Holger Hoos, Alan Hu, Margo Seltzer





Network Functions (NF) are useful and widespread

- Security
 - Firewall, DDoS protection, DPI
- Monitoring
 - QoE monitor, Network Stats
- Services
 - Ad insertion, Transcoder
- Network optimization
 - NAT, Load-balancer, WAN accelerator

























middleboxes ≈ # L2/L3 devices [Sherry et al. SIGCOMM'12]

Benefits of Virtualized Network Functions (VNF)

- Elasticity
 - Quick scale up and down NFs
- Fast upgrades
 - No need to wait for new hardware
- Quick configuration, recovery
 - Failover to the backup NF instance
- Outsourcing











firewall



QoE monitor



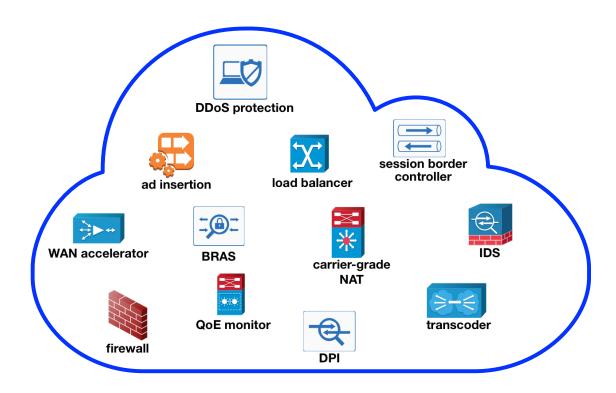




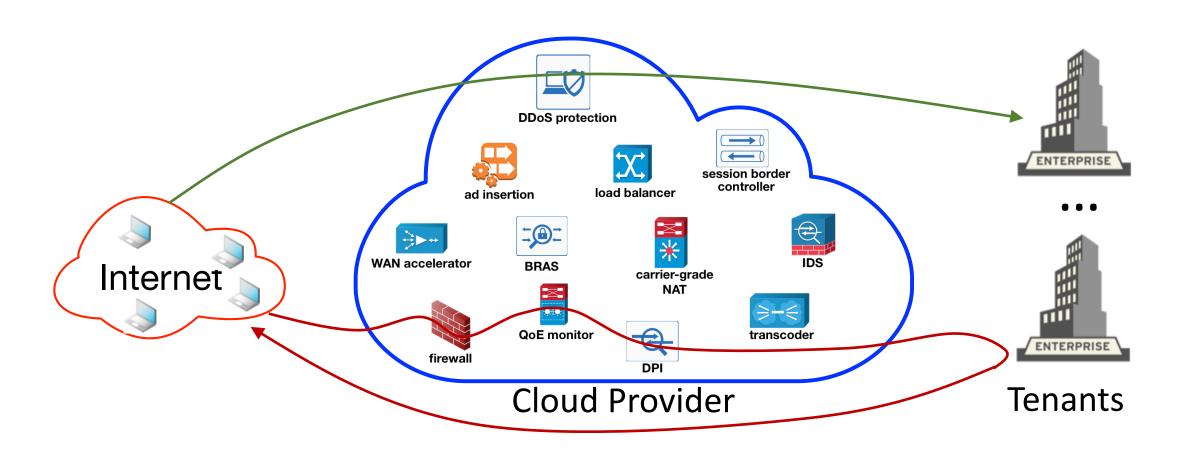




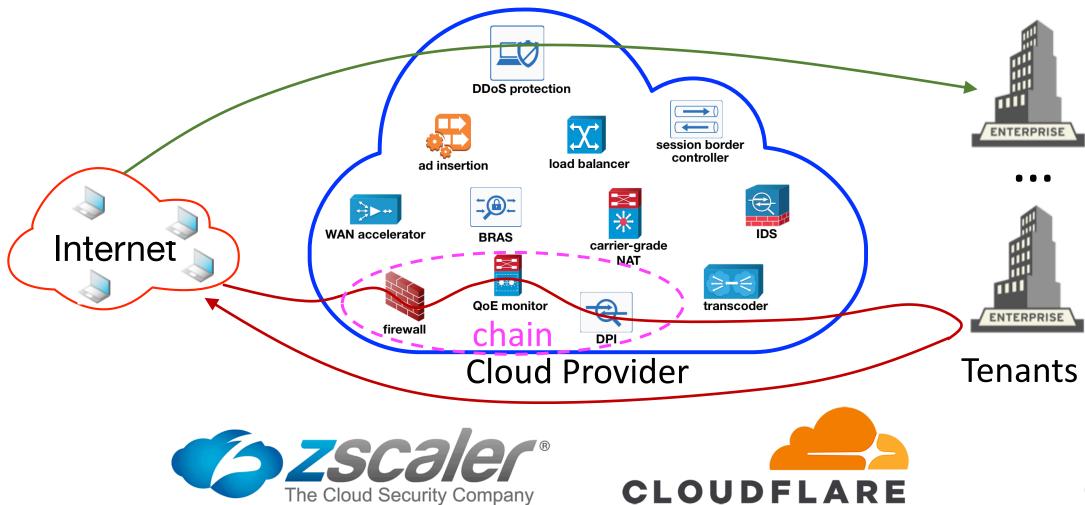
Outsourcing VNFs to the Cloud



Outsourcing VNFs to the Cloud



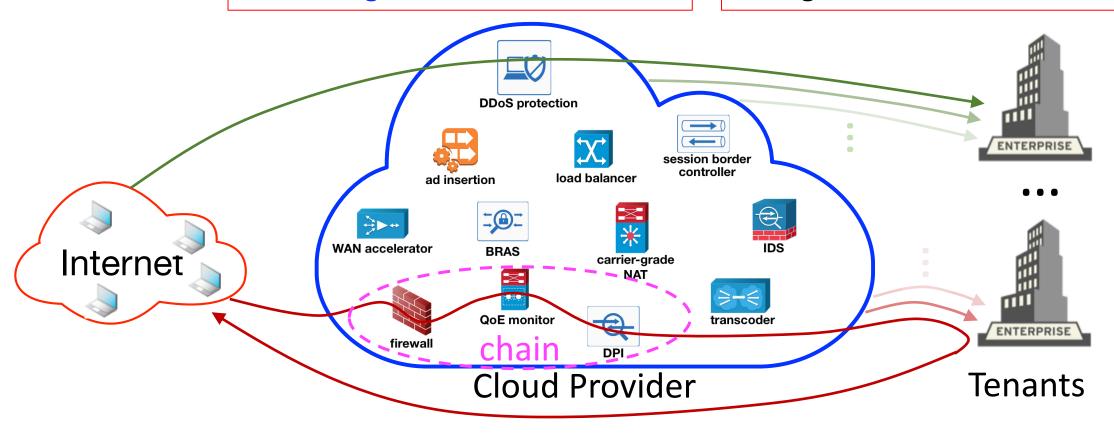
Outsourcing VNF Chains to the Cloud



Challenges of outsourcing VNF Chains

How can tenants allocate and manage their VNF chains?

How can cloud providers achieve high datacenter utilization?

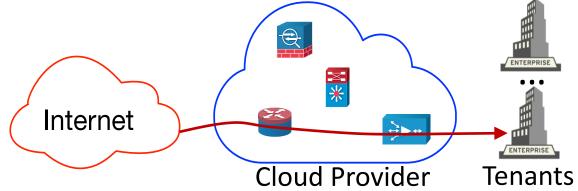


Our contributions: API and algorithm

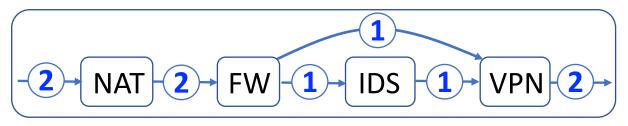
How can tenants allocate and manage their VNF chains?

How can cloud providers achieve high datacenter utilization?

- API to allocate and manage VNF chains
- Three algorithms
 - implement the API, and
 - achieve high datacenter utilization
- Evaluation
 - simulate: in datacenter scale with 1000+ servers
 - Daisy: emulate chain management at rack-scale
- Ongoing work: chain abstraction



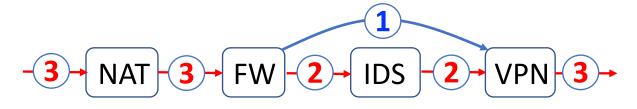
VNF Chain: six API with use-cases



Initial chain

cid ← allocate-chain(C, bw)
add-link-bandwidth(a, b, bw, cid)
add-node(f, cid)

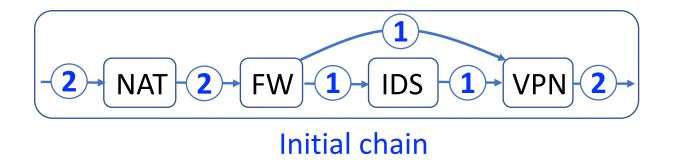
remove-link-bandwidth(a, b, bw, cid)
remove-node(f, cid)
remove-e2e-bandwidth(cid, bw)



NAT 2 FW 1 IDS 1 VPN 2 IDS' 1 Element upgrade

Chain scale-out

VNF Chain: six API with use-cases



A graph can be transformed arbitrarily by manipulating individual nodes and edges.

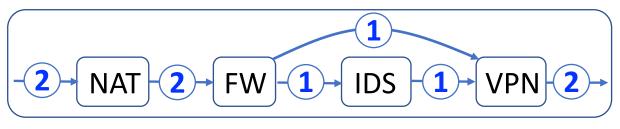
Chain scale-out

Element upgrade

Chain expand

• •

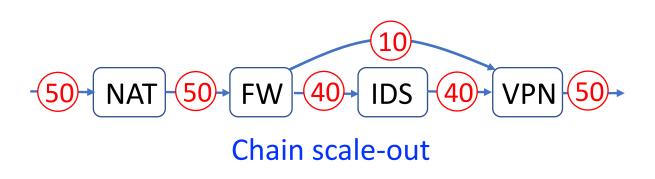
Scale-out beyond single physical resource capacity

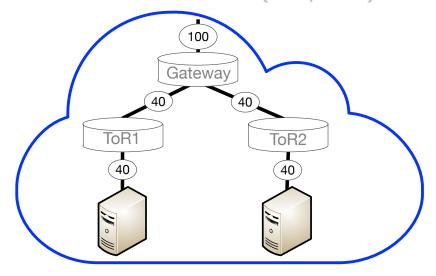


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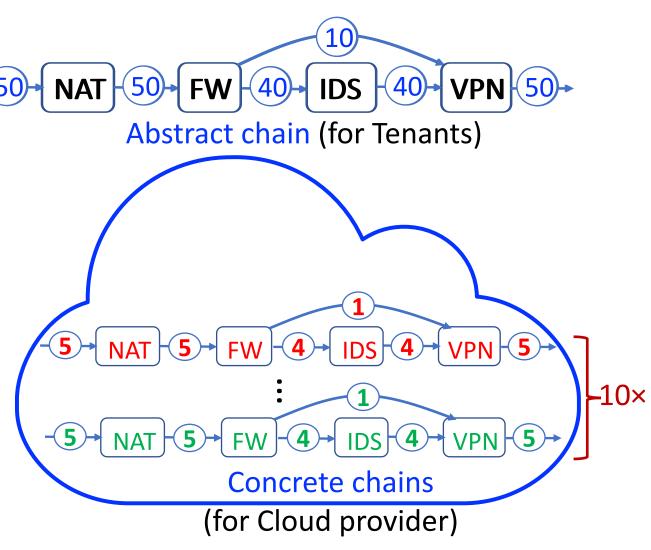
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Chain Abstraction: Abstract-Concrete VNF Chains

- Abstract VNF chain
 - what tenant requires to allocate and operates on
- Concrete VNF chain
 - cloud provider's implementation of the abstract chain
- Chains abstraction advantages
 - facilitates high DC utilization
 - improves SLA guarantees

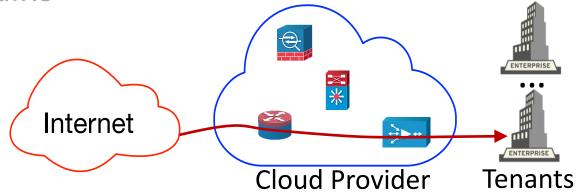


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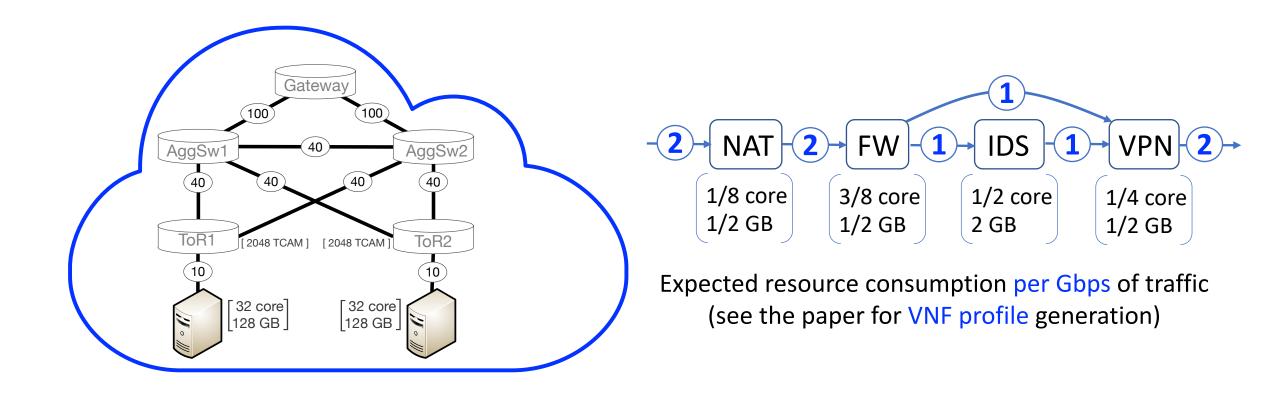
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How can tenants allocate and manage their VNF chains?

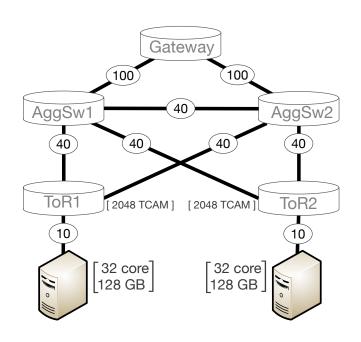
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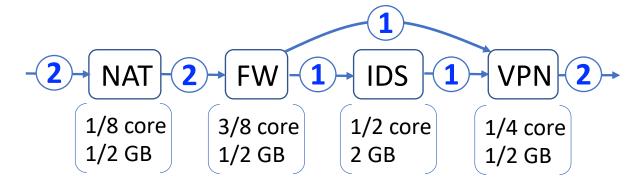


Algorithm inputs: DC topology and chain



Algorithms for Chain Allocation and Management

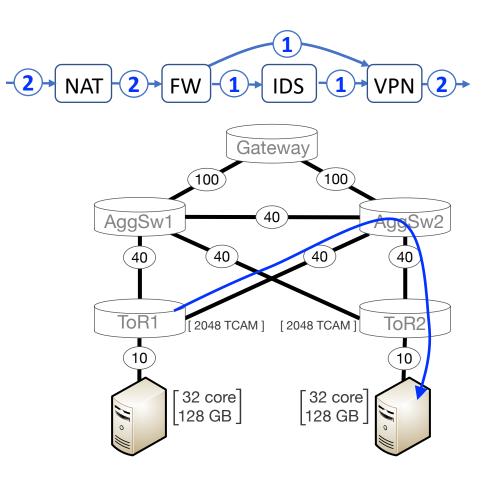




Expected resource consumption per Gbps of traffic (see the paper for VNF profile generation)

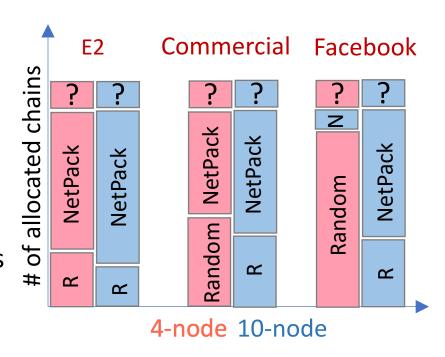
Algorithms for Chain Allocation and Management

- Random (baseline)
 - Consider NFs and servers/switches in random order
 - Attempt the above step n times (e.g., n=100)
 - Choose the shortest path between chain NFs



Algorithms for Chain Allocation and Management

- Random (baseline)
 - Consider NFs and servers/switches in random order
 - Attempt the above step n times (e.g., n=100)
 - Choose the shortest path between chain NFs
- NetPack: Random + 3 simple heuristics
 - Consider the chain NFs in a topological order
 - Re-use the same server when allocating consecutive NFs
 - Gradually increase the network scope: rack, cluster, etc.
- VNFSolver: how optimal is NetPack?
 - Constraint-solver based chain allocation algorithm
 - Slow, but complete: finds a solution when one exists

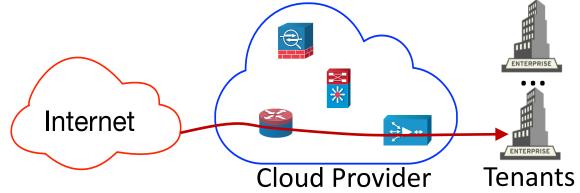


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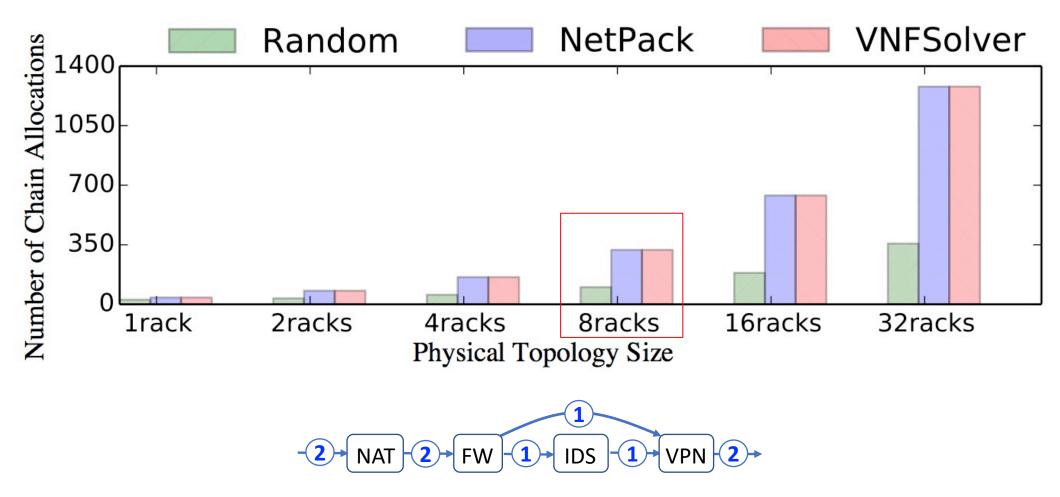
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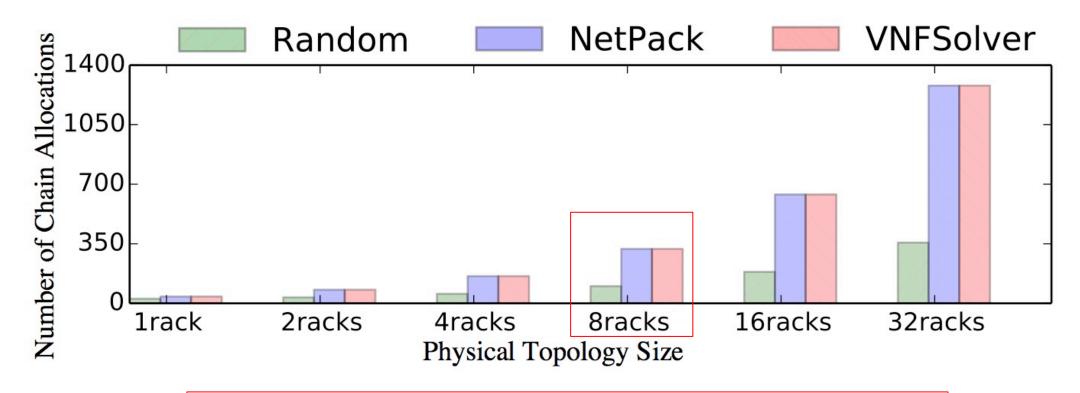
Evaluation: Objectives

- How good is the datacenter utilization?
 - Evaluate Random, NetPack, and VNFSolver
 - Consider three different datacenter topologies
 - Use five different VNF chains with varying length (2-10)
- How fast is chain allocation?
 - Measure time it takes to saturate the datacenter
- Does API reliably implement the use-cases?
 - Prototype scale-out and chain upgrade in Daisy
 - Use two different racks, two sources of packet traces

Datacenter utilization evaluation



Datacenter utilization evaluation



NetPack achieves at least 96% of VNFSolver allocations.

Chain allocation time: Random ≤ NetPack ≪ VNFSolver.

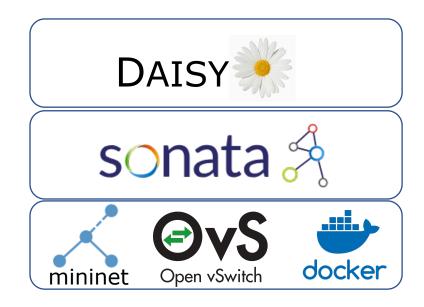
NetPack Utilization and Speed

NetPack achieves at least 96% of VNFSolver allocations while being 82x faster than VNFSolver (optimal) and only up to 54% slower (per chain) than Random (baseline) on average.

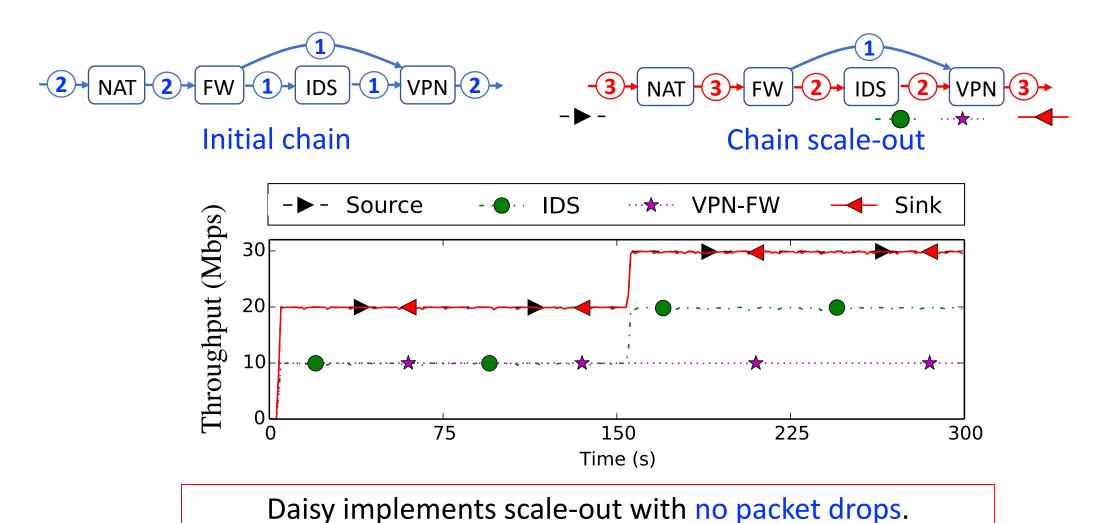
Qualitatively similar results with Facebook and Commercial DC topologies with chains of up to 10 nodes. (see the paper for details)

Feasibility check: does API work?

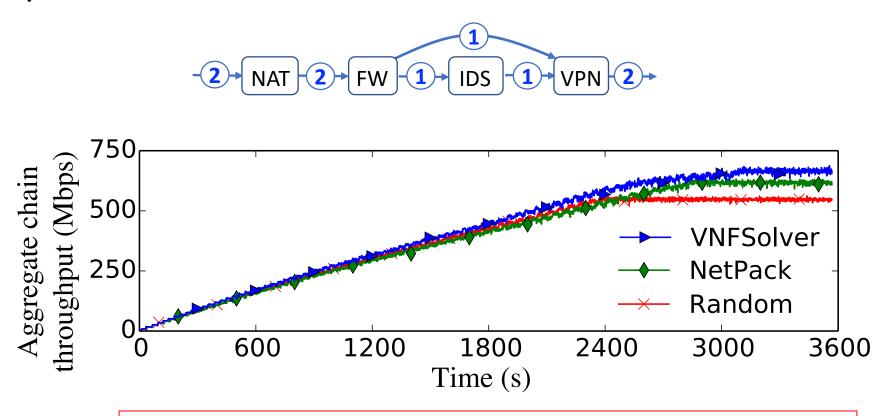
- Daisy builds on Sonata framework
 - Mininet to build DC topology
 - OVS for switches, and Dockers for NFs
- Runs on a single Azure VM
 - 64 cores, 432 GB RAM
- Emulates use-cases and chain arrivals
 - scale-out and upgrade use-cases
 - continuous arrival of tenant chains in rack-scale



VNF Chain scale-out



Daisy: emulate continuous chain arrival



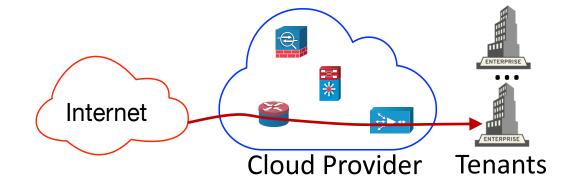
VNFSolver allocated 75 chains (687 Mbps)
NetPack allocated 67 chains (633 Mbps)
Random allocated 61 chains (561 Mbps)
(throughput with iperf generated packets is precise)

Daisy Contributions

Daisy implements scale-out with no packet drops and element upgrade with 1s packet drop at most.

We also emulated continuous chain arrival case where different tenants make chain allocation requests one-by-one.

Snapshot of complete work



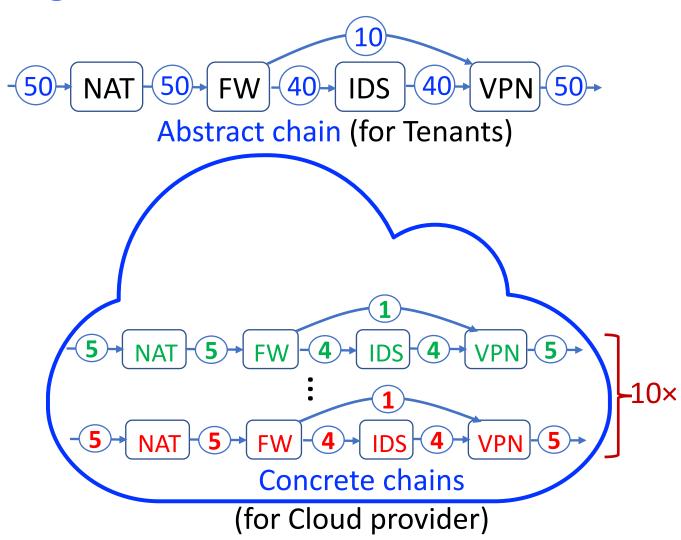
- API with six primitives
 - Implements wide-range of chain operations
 - Chain abstraction facilitates full DC utilization
- NetPack algorithm
 - Handles DC-scale allocation with 1000+ servers
 - Achieves at least 96% allocations of VNFSolver (optimal) while being 82x faster on average
- Daisy prototype
 - Demonstrates feasibility of API and algorithms
- Ongoing work: chain abstraction

How can tenants allocate and manage their VNF chains?

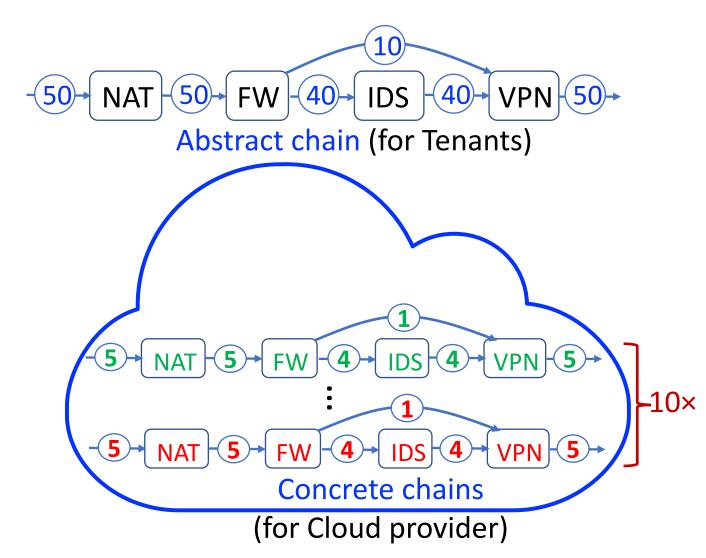
How can cloud providers achieve high data center utilization?

Chain abstraction challenges

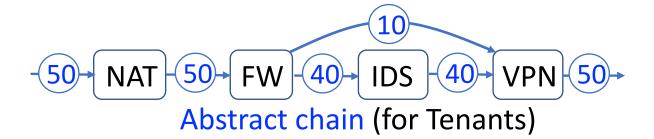
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 - what tenant requires to allocate and operates on
- Concrete VNF chain
 - cloud provider's implementation of the abstract chain
- Chains abstraction advantages
 - facilitates high DC utilization
 - improves SLA guarantees
- Challenges
 - low-latency, packet loss, state synchronization, efficiency loss, hotspots

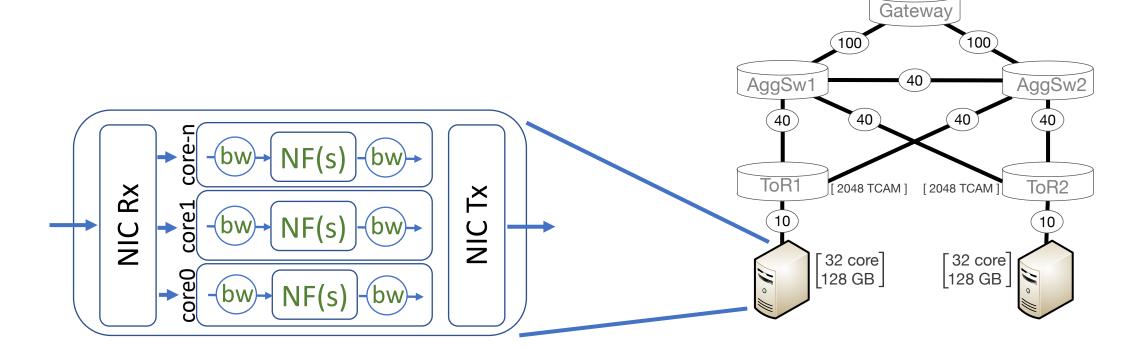


Problem: Hotspots



Problem: Hotspots

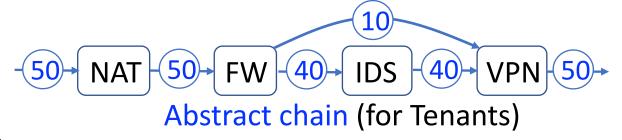


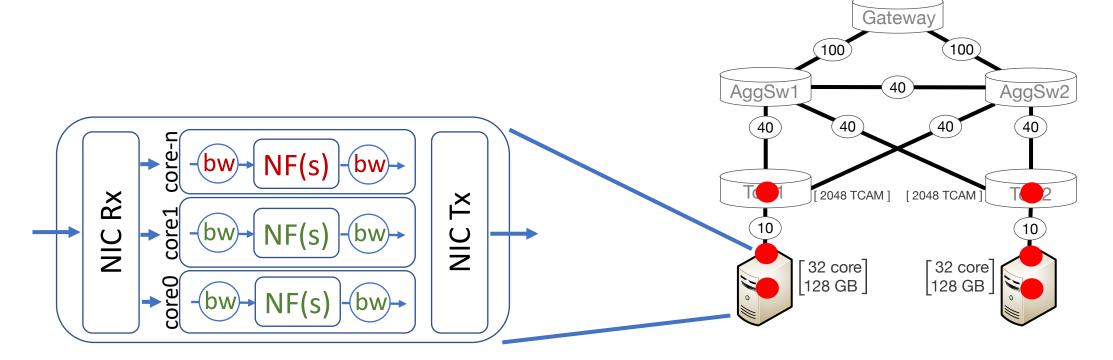


Problem: Hotspots

Hotspots in different layers:

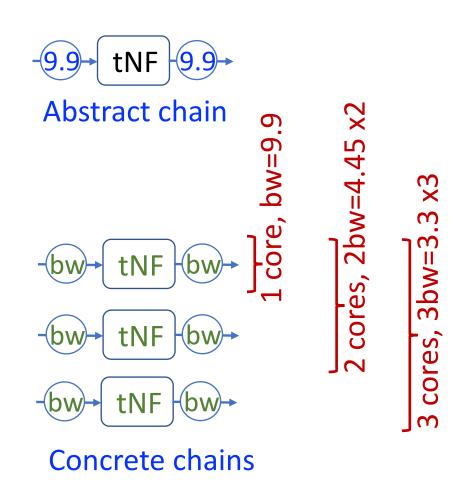
CPU cores, NIC ports, ToR switch ports, etc.





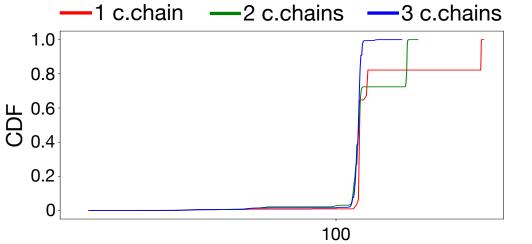
Microbenchmarks with up to 3 chains

- tNF: tunable NF
 - computes N prime numbers, per packet (<50K in our experiments)
 - Throughput: 0.83 Mpps (9900 Mbps with 1500 byte packet)
- NF runs on OpenNetVM
- Varied the number of concrete chains from one to three
- Each concrete chain processes a separate flow (5-tuple)



Microbenchmarks with up to 3 chains

- Tail latencies (99 percentile) in microseconds
 - 1 concrete chain (uses 1 core for 9.9Gbps): 336 us
 - 2 concrete chains (uses 2 cores for 9.9Gbps): 182 us
 - 3 concrete chains (uses 3 cores for 9.9Gbps): 126 us

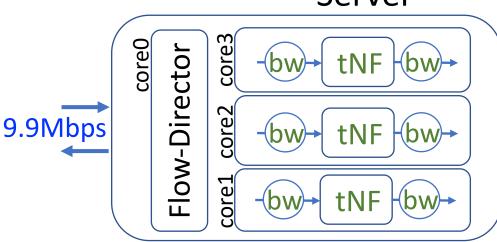


Packet processing latency in log scale (microseconds)



Abstract chain

Server



Concrete chains

Observations and Ongoing Work

- Latency grows proportional to the core load
- Need CPU load-aware chain-splitting mechanism
 - N flows to 1 core (N-to-1) for mice flows
 - 1-to-N for elephant flows
- Splitting should happen in multiple levels
 - CPU cores, NIC ports, ToR ports
- Need to support wide-range of NFs
 - Bounded tail latencies are particularly challenging for stateful NFs, such as DPI



Abstract chain







Concrete chains

Conclusion

Internet Cloud Provider Tenants

- API with six primitives
 - Implements wide-range of chain operations
 - Chain abstraction facilitates full DC utilization
- NetPack algorithm
 - Handles DC-scale allocation with 1000+ servers
 - Achieves at least 96% allocations of VNFSolver (optimal) while being 82x faster on average
- Ongoing work: chain abstraction
 - Need load-aware chain-splitting mechanism

How can tenants allocate and manage their VNF chains?

How can cloud providers achieve high data center utilization?

Thank you!